

IMAGE PROCESSING APPARATUS

BACKGROUND OF THE INVENTION

Field of the Invention

5 The present invention relates to an image processing apparatus for preventing the forgery of bills, etc. in a recording system for processing an image in a host computer and printing the image on a recording device.

10 Related Background Art

 Many systems for preventing the forgery of bills, valuable securities, etc. are designed as input/output device built-in systems such as copying units, etc. However, with an increasing number of personal
15 computers, high-performance peripheral devices such as a scanner, a digital camera, a printer, etc. have been developed. As a result, with these high-performance peripheral devices, the personal computers can output an image at a quality level higher than an image output
20 by the input/output device built-in copying units. Thus, in a personal computer environment, a forgery preventing system is required.

 As a characteristic of a forgery preventing system in the personal computer environment, a host computer
25 controls input and output devices. Therefore, it is necessary to allow an image process program of the host computer to recognize specified patterns of bills, valuable securities, etc. A personal computer can

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output data without a specific image process program only if it is informed of the control code of the output device.

Accordingly, a method of encrypting the control code of the output device is effective to prevent the use of a program other than the specific image process program for a forgery preventing process. The above mentioned method is disclosed by, for example, by Japanese Patent Application Laid-Open No. 6-105141 although not in the personal computer environment. FIG. 2 shows an example when this method is applied to the personal computer environment.

In FIG. 2, the processes in steps S2001 to S2005 are performed by the host computer, and the processes in steps S2006 to S2009 are performed by the recording device.

In the first step S2001, image data is input from an OS (basic software) or an application. Then, in step S2002, an image process including a forgery preventing process is performed. The image process in step S2002 includes color matching, gamma correction, and quantizing processes to convert an input image into print image data. In the forgery preventing process in step S2002, determination is normally made on a specified image by performing a pattern recognizing process. If the input image matches the specified image, then an error process is performed on the image

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data.
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5 data as a command for control of a printer in step
S2004.

10 circuit (not shown in the attached drawings) in step
S2005.

15 S2007, and encrypted print image data is generated.

medium in step S2009.

20 However, since in the above mentioned system which
performs the encryption process between the host
computer and the printer, there is an one to one
correspondence between the input image and the data
transferred to the printer, the encrypted data can be
25 easily decrypted.

control code of the printer used in the conventional

system is not open, and there is a problem that the image forgery preventing process cannot be effectively performed.

5 SUMMARY OF THE INVENTION

An object of the present invention is to solve the above mentioned problems, and provide an improved image processing apparatus and method.

10 Another object of the present invention is to provide an image processing apparatus and method capable of effectively performing a forgery preventing process on an input image.

15 A further object of the present invention is to provide an image processing apparatus and method by hardening the decryption of a forgery preventing process.

20 A further object of the present invention is to provide an image processing apparatus and method capable of effectively preventing a print process performed by a number of unspecified image recording devices by encrypting print control data for controlling an image recording device using a common key issued by an image recording device when the image recording device is allowed to print an image, and by
25 disabling any other image recording devices than the image recording device by which the common key has been generated to record the image.

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Further objects of the present invention will be clearly described by the following explanation based on the attached drawings and the claims.

5 BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a flowchart of the procedure of the operations of the recording system according to the first embodiment of the present invention;

FIG. 2 is a flowchart of the procedure of the encryption process according to the recording system of the conventional technology;

FIG. 3 is a flowchart of the procedure of an encryption process according to the first embodiment of the present invention;

FIG. 4 is a flowchart of the procedure of a decryption process according to the first embodiment of the present invention;

FIG. 5 is a flowchart of the procedure of an encryption table generating process according to the first embodiment of the present invention;

FIG. 6 is a flowchart of the procedure of a decryption table generating process according to the first embodiment of the present invention; and

FIG. 7 is a block diagram of an example of the configuration of the image data recording system realizing the present invention.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS
(First Embodiment)

5 The embodiments of the present invention are described below in detail by referring to the attached drawings.

FIG. 1 is a flowchart of the procedure of processing data in the recording system according to the present invention.

10 In FIG. 1, the processes in steps S1001 to S1004 and steps S1009 to S1012 are performed by a host computer. The processes in steps S1005 to S1008 and steps S1013 to S1017 are performed by a recording device.

15 In the first step S1001, image data is input. Then, in step S1002, an image process including a forgery preventing process is performed.

Next, in step S1003, a print ID is generated, and then the print ID is transferred to a recording device in step S1004.

20 Next, in step S1005, the recording device receives a print ID, and stores the received print ID in step S1006.

25 Next, in step S1007, a common key is generated. At this time, the recording device is managed such that a print ID is paired with a common key.

If a common key is generated in step S1007, it is irregularly generated such that the common key cannot

be associated with the value of the print ID.

Then, in step S1008, the generated common key is transmitted to the host computer.

5 The common key transmitted from the recording device is received by the host computer in step S1009. Thus, since the recording device issues the common key in response to the transfer of the print ID of the host computer, the transfer of the print ID in step S1004 indicates a request to issue a common key.

10 In step S1010, the print image data generated in the process performed in step S1002 is encrypted by the common key received in step S1009.

15 Then, in step S1011, the encrypted print image data is converted into a print control command. Then, in step S1012, the print ID and the print control data command are transferred to the recording device side.

20 In step S1013, the recording device receives a print ID and print control data. Then, in step S1014, the common key corresponding to the received print ID is retrieved and obtained from a managed pair of the print ID and the common key. Then, in step S1015, the command of the print control data is analyzed, and the encrypted print image data is extracted.

25 Then, in step S1016, the print image data is decrypted using the common key obtained in step S1014. In step S1017, the print image data is stored on the storage medium in the print process.

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The print data transferring process in steps S1012 and S1013 and the process in steps S1015 to S1017 can be sequentially repeated in parallel. Although not shown in FIG. 1, the used print ID and the common key
5 paired therewith may be discarded after performing the print process in step S1017.

FIG. 3 is a flowchart of the contents of the encryption process performed in step S1010.

The random number table according to this
10 embodiment comprises series of irregularly arranged integers (1-byte length each) of "0" to "255". That is, the table size is 256 bytes. The common key is formed by integers from 0 to 255.

First, in step S3001, a random number table is
15 stored on the memory (RAM) of the host computer. Then, in step S3002, the random number table on the memory (RAM) is converted into an encryption table using a common key.

In step S3003, the print image data is encrypted
20 using the generated encryption table. In this case, the print image data is sequentially read in a byte unit, the data value of one read byte is set as an offset from the leading address of the encryption table, and the value of the corresponding address is
25 set as the encryption print image data.

FIG. 5 is a flowchart of the process performed in step S3002.

FIG. 3 is a flowchart of the contents of the encryption process performed in step S1010.

First, in step S5000, the process is started.
then, in step S5001, a variable n is set to "0". The
variable n is a management counter for 100 times
repetition of the processes in steps S5003 to S5005
5 described later.

Then, in step S5002, the value of the common key
is assigned to a variable B. Then, in step S5003, a
variable A is computed using the variable B by the
following Equation (1). 'mod' is a well-known function
10 for use in obtaining a residue in a division.

$$A = (5 \times B + 13) \bmod 256 \dots (1)$$

Next, in step S5004, the variable B is computed
using the variable A by the following Equation (2).

$$B = (5 \times A + 13) \bmod 256 \dots (2)$$

15 In the computation by the Equations (1) and (2),
pseudo random numbers are generated in a linear
congruential method. That is, the common key is used
for the initial value in the linear congruential
method.

20 Then, in step S5005, the table value whose offset
from the leading address of the random number table
stored in the memory in step S3001 is A and the table
value whose offset is B are interchanged. Then, in
step S5006, n is incremented by 1.

25 It is determined in step S5007 whether or not n is
"100". If it is "100", then control is passed to step
S5008, and the converting operation terminates. If it

is determined in step S5007 that n is not 100, control is passed to step S5003.

FIG. 4 is a flowchart explaining the procedure of the decryption process performed in step S1016. The random number table used in this embodiment is the same table as that used in the process in step S3001.

First, in step S4001, the random number table is developed on the memory (RAM) of the recording device.

Then, in step S4002, the random number table on the memory (RAM) is converted into a decryption table using a common key. In the obtained decryption table, the offset value from the leading address of the series and the integer stored at the address can be obtained by interchanging the values on the above mentioned encryption table.

For example, when the 25th value from the start of the encryption table is "12", the 12th value from the start of the decryption table is "25" (assume that the start of the table is set to 0).

That is, the encryption table is an inverse conversion table of the decryption table. Assuming that a function A() indicates the conversion using the encryption table and a function B() indicates the conversion using the decryption table, the following Equations exist.

$$a = A(d), d = B(a)$$

Then, the encryption print image data is decrypted

using the generated decryption table in step S4003. In
step S4003, the encryption print image data is
sequentially read in a byte unit. The read 1-byte data
value is an offset from the leading address of the
5 encryption table, and the value of the corresponding
address is the print image data.

FIG. 6 is a flowchart of the process in step
S4002.

10 In FIG. 6, the processes in steps S6000 to S6007
are the same as the converting operation in steps S5000
to S6007 on the encryption table. After generating the
encryption table, the address value and the table value
are interchanged in step S6008 for conversion into the
decryption table.

15 The following Tables 1, 2, and 3 are the tables
used or generated in this embodiment.

TOP SECRET

Random Number Table

x	R(x)
0	88
1	197
2	230
3	139
4	196
5	225
6	114
7	71
8	112
9	61
10	62
11	67
12	92
13	217
14	74
15	127
16	136
17	181
18	150
19	251
20	244
21	209
22	34
23	183
24	160
25	45
26	238
27	179
28	140
29	201
30	250
31	239
x	R(x)
32	184
33	165
34	70
35	107
36	36
37	183
38	210
39	39
40	208
41	29
42	158
43	35
44	188
45	185
46	170
47	95
48	232
49	149
50	246
51	219
52	84
53	177
54	130
55	151
56	0
57	13
58	78
59	147
60	236
61	169
62	90
63	207
x	R(x)
64	24
65	133
66	166
67	75
68	132
69	161
70	50
71	7
72	48
73	253
74	254
75	3
76	28
77	153
78	10
79	63
80	72
81	117
82	86
83	187
84	180
85	145
86	226
87	119
88	96
89	237
90	174
91	115
92	76
93	137
94	186
95	175
x	R(x)
96	120
97	101
98	6
99	43
100	228
101	129
102	146
103	231
104	144
105	221
106	84
107	227
108	124
109	121
110	106
111	31
112	168
113	85
114	182
115	155
116	20
117	113
118	66
119	87
120	192
121	205
122	14
123	83
124	172
125	105
126	26
127	143
x	R(x)
128	216
129	69
130	102
131	11
132	68
133	97
134	242
135	199
136	240
137	189
138	190
139	195
140	220
141	89
142	202
143	255
144	8
145	53
146	22
147	123
148	116
149	81
150	162
151	55
152	32
153	173
154	110
155	51
156	12
157	73
158	122
159	111
x	R(x)
160	56
161	37
162	198
163	235
164	164
165	65
166	82
167	167
168	80
169	157
170	30
171	163
172	60
173	57
174	42
175	223
176	104
177	21
178	118
179	91
180	212
181	49
182	2
183	23
184	128
185	141
186	206
187	19
188	108
189	41
190	218
191	79
x	R(x)
192	152
193	5
194	38
195	203
196	4
197	3

Table 3

Decryption Table

x	B(x)
0	37
1	248
2	155
3	166
4	196
5	36
6	247
7	71
8	221
9	144
10	147
11	126
12	25
13	60
14	111
15	202
16	149
17	213
18	230
19	86
20	81
21	84
22	231
23	183
24	64
25	192
26	131
27	46
28	137
29	108
30	95
31	122

x	B(x)
32	5
33	197
34	123
35	6
36	193
37	132
38	194
39	210
40	189
41	240
42	115
43	99
44	252
45	156
46	79
47	42
48	117
49	181
50	107
51	182
52	49
53	180
54	199
55	130
56	160
57	32
58	222
59	211
60	105
61	204
62	63
63	218

x	B(x)
64	229
65	184
66	91
67	102
68	161
69	228
70	34
71	50
72	157
73	80
74	83
75	67
76	92
77	249
78	47
79	138
80	85
81	232
82	75
83	22
84	17
85	20
86	167
87	226
88	13
89	128
90	62
91	238
92	73
93	233
94	31
95	58

x	B(x)
96	88
97	24
98	59
99	198
100	129
101	68
102	151
103	146
104	125
105	176
106	51
107	158
108	185
109	217
110	154
111	159
112	53
113	72
114	43
115	118
116	148
117	116
118	135
119	66
120	237
121	224
122	35
123	78
124	41
125	140
126	255
127	15

x	B(x)
128	165
129	120
130	54
131	38
132	97
133	164
134	119
135	242
136	16
137	93
138	19
139	254
140	153
141	188
142	239
143	127
144	21
145	168
146	11
147	214
148	244
149	212
150	18
151	162
152	205
153	77
154	3
155	174
156	9
157	236
158	223
159	250

x	B(x)
160	133
161	216
162	150
163	171
164	65
165	4
166	87
167	82
168	112
169	61
170	243
171	94
172	121
173	28
174	207
175	170
176	245
177	8
178	235
179	27
180	177
181	52
182	114
183	2
184	173
185	45
186	227
187	14
188	44
189	76
190	191
191	90

x	B(x)
192	101
193	56
194	219
195	139

The above mentioned Table 1 is an example of the random number table according to this embodiment. The Table 2 is an encryption table when the common key is "15" in this embodiment. The Table 3 is a decryption table when the common key is "15" in this embodiment.

In the above mentioned embodiment, the encryption process and the decryption process are performed by the conversion system using the tables. This system requires a smaller load in performing an operation, thereby preventing the printing speed from being reduced.

Furthermore, according to this embodiment, the recording device manages both print ID and common key. Therefore, when the recording device is connected to a plurality of host computers, the print process can be performed in the transfer order of print control data regardless of the common key issue order.

(Other Embodiments)

In the above mentioned embodiments, an encryption table is generated by converting a random number table using a common key, but the encryption table can be the common key. Furthermore, the common key can be used as a parameter in the congruential method, and an encryption table can be generated by generating pseudo random numbers.

The encryption table may be generated by conversion using the common key according to not only

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the congruential method but also an average method.

Since a common key of a specified value can be issued in the above mentioned embodiments, the value of an internal timer of the recording device can be used.

5 An example of the configuration of the image data recording system embodying the present invention is described below by referring to the block diagram shown in FIG. 7.

10 In FIG. 7, reference numeral 70 denotes an image data processing device; reference numeral 71 denotes an interface; reference numeral 72 denotes an image processing unit; reference numeral 73 denotes a print ID generating unit; reference numeral 74 denotes a print ID storage unit; reference numeral 75 denotes a
15 first transfer unit; reference numeral 76 denotes an encryption unit; reference numeral 77 denotes a print control data generating unit; and reference numeral 78 denotes a second transfer unit.

20 Furthermore, reference numeral 80 denotes an image data recording device; reference numeral 81 denotes an interface; reference numeral 82 denotes a common key generating unit; reference numeral 83 denotes a management unit; reference numeral 84 denotes a common key issue unit; reference numeral 85 denotes a common
25 key obtaining unit; reference numeral 86 denotes an analyzing unit; reference numeral 87 denotes a decryption unit; and reference numeral 88 denotes a

print unit.

As shown in FIG. 7, the image data recording system comprises the image data processing device 70 and the image data recording device 80. Through the interface 71 and the interface 81 respectively provided in the image data processing device 70 and the image data recording device 80, various data and commands are transmitted and received to print the image data input to the image data processing device 70 on the image data recording device 80, and then output.

In FIG. 7, the image processing unit 72 performs the image process including the forgery preventing process on the input image data. The print ID generating unit 73 generates a print ID corresponding to the image data processed by the image processing unit 72 in the forgery preventing process. The generated print ID is stored in the print ID storage unit 74, and transferred by the first transfer unit 75 to the image data recording device 80.

The encryption unit 76 encrypts the image data processed by the image processing unit 72 in the predetermined processes using the common key transmitted from the image data recording device 80. The print control data generating unit 77 generates print control data by converting the print image data encrypted by the encryption unit 76 into a print control command. Then, the generated print control

data and the print ID generated by the print ID generating unit 73 and stored in the print ID storage unit 74 are transferred by the second transfer unit 78 to the image data recording device 80.

5 The common key generating unit 82 generates a common key based on the print ID transferred from the image data processing device 70. Then, the generated common key and the transferred print ID are stored and managed in the memory of the management unit 83.

10 The common key issue unit 84 transmits the common key generated by the common key generating unit 82 to the image data processing device 70. When the print ID and the print control data are transmitted from the image data processing device 70, the common key
15 obtaining unit 85 obtains a common key corresponding to the print ID from the management unit 83.

 The analyzing unit 86 analyzes the command of the above mentioned print data using the common key obtained by the common key obtaining unit 85, and
20 extracts the encrypted print image data. The decryption unit 87 decrypts the print image data extracted by the analyzing unit 86 using the common key obtained by the common key obtaining unit 85. Then, the decrypted print image data is stored in a storage
25 medium (not shown in the attached drawings) by the print unit 88.

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(Other Embodiments according to the Present Invention)

The present invention can be applied to either a system comprising a plurality of appliances (for example, a host computer, an interface appliance, a leader, a printer, etc.) or a device comprising one appliance.

In addition, the present invention further includes an embodiment in which a program code of software for realizing the functions of the above mentioned embodiments are provided for the computer in the device or the system connected to each of the above mentioned devices so that various devices can be operated to realize the functions of the above mentioned embodiments, and the above mentioned devices are operated according to the program stored in the computer (CPU or MPU) of the system or the device.

In this case, the program code of the software realizes the functions of the above mentioned embodiment, and the program code itself and the unit for providing the program code for the computer, for example, a storage medium storing the program code constitute the present invention. The storage medium storing the program code can be, for example, a floppy disk, a hard disk, an optical disk, a magneto-optical disk, CD-ROM, a magnetic tape, a non-volatile memory card, ROM, etc.

In addition, it is needless to say that the

program code is included in the embodiment of the present invention not only in the case where the function described in the above description of the embodiment is realized by a computer executing the program code provided thereto, but also in the case the function is realized by the program code cooperating with the OS (operating system) or other application software operating in the computer.

Furthermore, after the provided program code is stored in the memory in a function extension board of a computer or a function extension unit connected to the computer, the CPU, etc. provided in the function extension board or the function extension unit can perform all or part of an actual process at an instruction of the program code to realize the function of the above mentioned embodiment in the process. The present invention can also include the above mentioned case.

As described above, the print control data is encrypted using a common key issued by an image data recording device. Therefore, when the image data processing device generates a recording image by controlling the image data recording device, the image data recording device which has issued the common key is required, thereby effectively preventing a number of unspecified image data recording devices from printing image data, and ensuring that the forgery preventing

process is performed by the image data processing device. Thus, the forgery of bills, valuable securities, etc. can be prevented without fail.

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